Discrete Sampling-Based Planning

Dapeng Zhao (Eagle) dapengz Rockey Hester (Quint) rockeyh

Viraj Parimi vparimi

December 20, 2020

1 Abstract

We introduce dRRT, a discrete sampling-based planner that uses Bresenham's line algorithm and techniques from search-based planning to achieve anytime qualities and reuse computation. Results on planning for a 5-DOF robotic arm demonstrate faster performance than traditional RRT.

2 Introduction

For the project, we devised a discrete sampling-based planner called dRRT. Conventional samplingbased planners like RRT execute in the continuous configuration space (C-space). The problem with such approaches is that since their underlying sample space is so large, one needs to run them for multiple iterations to generate a plan, even in trivial cases. Variants of RRT like RRT-Connect and RRT* have been proposed to enhance convergence time and solution quality, respectively.

As an alternative, we propose a discrete version of RRT by leveraging ideas from search-based planning algorithms. Intuitively, using this approach, one would be able to generate the tree more efficiently. Further, one can even improve the collision checker's performance by utilizing the hierarchical decomposition of the problem.

3 Approach

3.1 Planning Representation

In order to represent the problem, we utilize the following planning representation,

$$M^R = < \theta >$$

 $M^W = < \text{obstacle/free space} >$
 $S^R_{current} = < \theta_{current} >$
 $S^W_{current} = < \text{constant} >$
 $C = < \text{Euclidean Distance} >$
 $G = < \theta_{acal} >$

where M^R represents the model of the robot, M^W represents the model of the world, $S^R_{current}$ represents the current model of the robot, $S^W_{current}$ represents the current model of the world, C represents the cost function, and G represents the goal. For experiments, we utilize the setup provided in HW2 and consequently implement the planner on a 5-DOF arm robot. This setup implies that θ is a 5-dimensional variable of joint angles.

3.2 Algorithm

Algorithm 1 presents the pseudo-code of our proposed algorithm. The algorithm's input is the q_{init} and q_{goal} configurations of the robot. Since these configurations lie in the continuous space, we first utilize the LocateOnGrid function to map them to particular points on the grid representation of the C-space. To generate this grid, we start with a resolution λ of 2 and keep increasing this in multiples of 2 until we can map q_{init} and q_{goal} to grid points. Once mapped, we add the grid point v_{init} to the Tree \mathcal{T} .

We then sample a random configuration q_{rand} from the grid C-space by utilizing the Sample function. To mitigate situations where it's hard to find the goal, we utilize goal-bias sampling. This function's key idea is that we maintain a set \mathbb{Q} of *not-to-draw points*. This set consists of grid points that have been either sampled or added to the Tree \mathcal{T} and is analogous to the *closed list* in search-based planning. Once we get q_{rand} we use the Extend function where like the regular RRT we detect the nearest grid point v_{near} that is part of \mathcal{T} from q_{rand} . To identify which grid points would allow to reach q_{rand} we exploit the Bresenham function. For each of the points v_{new} returned by it, we check whether the line connecting them is valid and is collision-free. If it was then, we add v_{new} to both \mathcal{T} and \mathbb{Q} .

When all the possible grid points at a given resolution λ get exhausted either because they were added to \mathcal{T} or because they were in collision, we increase the resolution of the underlying grid by using the Inflate function. We carefully update the nodes in \mathcal{T} and \mathbb{Q} except v_{init} and update their corresponding parents within this function. This technique allows us to skip the collision checking in finer resolutions for already existing nodes, thereby saving a lot of time even though the grid size is huge. This idea was also the reason as to why we chose to connect grid points greedily via **Bresenham** function because connecting two nodes without the intermediate points would violate this re-usability. Finally, when the algorithm can extend to q_{qoal} , we terminate and publish the path.

4 Experimental Analysis

4.1 Setup

There are ten pairs of start-goal, and each pair is evaluated **5** times for each algorithm, instead of only running once. If at one test, the planner couldn't find any viable path, the test is deemed as failed, and the number of this test is not considered for evaluation. Since the primary point of comparison is with RRT, we chose random start and goal configurations that favored RRT over the other algorithms. This way, we ensured that we could have a fair comparison between RRT and dRRT. A brief video showcasing our experiments can be found at https://youtu.be/Pwc9E1cjk7A. Our complete implementation can be found at https://bitbucket.org/eaglez1111/16782_final_proj/

4.2 Metrics Definition

The metrics used for evaluation:

- Time(s): How long the planner takes to return a feasible plan
- Succ.(%): Success rate
- #Spl.: The number of nodes in the tree/graph
- Cost: The euclidean length of the path in C-space
- Leng.: The number of configurations in the returned plan, the length of the plan

Algorithm 1: Discrete RRT

```
1 Function dRRT(q_{init}, q_{goal}):
             LocateOnGrid([q_{init}, q_{goal}])
  \mathbf{2}
             \mathcal{T}.\mathrm{add}(\mathbf{v}_{init} = \{q_{init}, \oslash\})
 3
             while q_{goal} \notin \mathcal{T} do
  4
                    q_{rand} \leftarrow \texttt{Sample}(\lambda)
  5
                    if q_{rand} then
  6
                          \texttt{Extend}(\mathcal{T}, q_{rand})
  7
                    else
  8
                           \texttt{Inflate}(\lambda)
  9
10 Function LocateOnGrid(Q):
             \lambda \leftarrow 2
11
             while \neg \text{Discretize}(\mathbb{Q}, \lambda) do
12
                    \lambda \leftarrow \lambda * 2
13
             return \lambda
\mathbf{14}
15 Function Extend(\mathcal{T}, q_{rand}):
             \mathbf{v}_{near} \leftarrow \texttt{NearestNeighbor}(q, \mathcal{T})
16
             \mathbb{Q}_{line} \leftarrow \texttt{Bresenham}(\mathbf{v}_{near}.q,q_{rand})
\mathbf{17}
             for q_{new} \in \mathbb{Q}_{line} do
\mathbf{18}
                    if IsValidLine(\mathbf{v}_{near}.q, q_{new}) then
19
\mathbf{20}
                            \mathbf{v}_{new}.q \leftarrow q_{new}
\mathbf{21}
                            \mathbf{v}_{new}.parent \leftarrow \mathbf{v}_{near}
                           \mathcal{T}.add(\mathbf{v}_{new})
\mathbf{22}
                            \overline{\mathbb{Q}}.add(\mathbf{v}_{new}.q)
23
                            \mathbf{v}_{near} \leftarrow \mathbf{v}_{new}
\mathbf{24}
                    else
\mathbf{25}
                           return NotReached
\mathbf{26}
             return Reached
\mathbf{27}
28 Function Sample(\lambda):
             if goal bias sampling then
\mathbf{29}
30
                    return q_{goal}
             else
\mathbf{31}
                    if \|\overline{\mathbb{Q}}\| == \lambda^D then
32
                        return \oslash;
33
                    \aleph_{\lambda} \leftarrow \text{all points at current resolution } \lambda
\mathbf{34}
                     q_{rand} \leftarrow (\mathbb{Q}_{pool} = \aleph_{\lambda} - \overline{\mathbb{Q}})[\texttt{RandomInt}()]
\mathbf{35}
                    \mathbb{Q}.add(q_{rand})
36
                    return
37
38 Function Inflate(\lambda):
             \lambda \leftarrow \lambda * 2
39
             for \mathbf{v}_i \in \mathcal{T} \setminus \{v_{init}\} do
40
                    \mathbf{v}_{new}.q \leftarrow (\mathbf{v}_i.q + \mathbf{v}_i.parent.q)/2
\mathbf{41}
                    \mathbf{v}_{new}.parent \leftarrow \mathbf{v}_i.parent
\mathbf{42}
                    \mathbf{v}_i.parent \leftarrow \mathbf{v}_{new}
43
                    \mathcal{T}'.add(\mathbf{v}_{new})
\mathbf{44}
                    \overline{\mathbb{Q}}.add(\mathbf{v}_{new}.q)
\mathbf{45}
             \mathcal{T}.append(\mathcal{T}')
46
```

4.3 Performance (mean)

Average values of the metrics are shown here:

Algo	Time(a)	Suce (%)	#Spl	Cost	Long
Algo	1 me(s)	Succ.(70)	#spi.	COSt	Leng.
RRT	0.076	100.0	255.8	7.95	12.50
dRRT	0.027	100.0	3.6	10.58	4.80
RRT-Connect	0.008	100.0	343.7	14.36	23.83
RRT-Star	0.018	100.0	242.9	8.03	12.90
PRM	0.770	83.3	3002.0	4.00	2.83

Table 1: Average values of the metrics

4.4 Performance (variance)

For each state-goal pair and each algorithm, the variance of the 5 trials is calculated; variances are then averaged for each algorithm to be reported here:

Algo	Time	Succ.	$\#$ Spl. (10^5)	Cost	Leng.
RRT	0.141	0.0	3.7143	16.33	52.92
dRRT	0.004	0.0	0.0001	0.09	3.47
RRT-Connect	0.000	0.0	0.2697	9.58	24.14
RRT-Star	0.000	0.0	0.3127	7.94	33.36
PRM	0.003	0.1	0.0000	0.73	0.13

Table 2: Variance values of the metrics

4.5 Discussion

- dRRT performed faster than RRT on average, which shows the performance boost discrete sampling provided mainly due to reusing of computation for collision checking. However, RRT-Connect and RRT* were better performing algorithms in terms of time. In terms of success rate, all the algorithms except PRM performed flawlessly. This situation happens due to the biased configurations which we took, as mentioned previously.
- The number of nodes within the tree is minimal for dRRT compared to others, which happens since we are working in the discrete C-space. This allows us to skip large distances over the C-space, thereby allowing us to connect far-off configurations. This inference is further corroborated by the number of nodes in the final path. dRRT and PRM returned smaller paths compared to other algorithms.
- The only metric where dRRT is beaten by RRT is the path cost. This finding was expected due to the huge skips we make over the C-space. One potential improvement of dRRT would be to optimize the path to return lower-cost solutions.
- dRRT algorithm provides resolution-based completeness guarantee where if the solution was to exist in a particular resolution, the algorithm is bound to find it. This is true because before we move to a finer resolution, we exhaust all the possible grid points in the current resolution. By doing that, we ensure that if there were a potential solution, dRRT would never miss it.
- The hierarchical decomposition routine allows us to propose a quick solution very fast at higher resolutions. This anytime solution can be further improved depending on how much time is available by moving down the hierarchy and working with finer resolutions.

• Low variance on performance metrics is essential. dRRT is more stable and predictable, but we believe that this could be because of very few nodes in the tree. For metrics like path cost and runtime, being consistent helps algorithm designers and programmers better manage the system.

5 Conclusion

We developed and evaluated dRRT, a new algorithm for discrete sampling-based planning that sacrifices initial solution quality in order to outperform RRT in time and space complexity. Future work entails development of discrete versions of other sampling-based planners such as RRT*, RRT-Connect, and PRM, as well as developing an efficient post-processing tool to smooth dRRT solutions.